

Practical DNSSEC Debugging

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With thanks to Olaf Kolkman & Willem Toorop

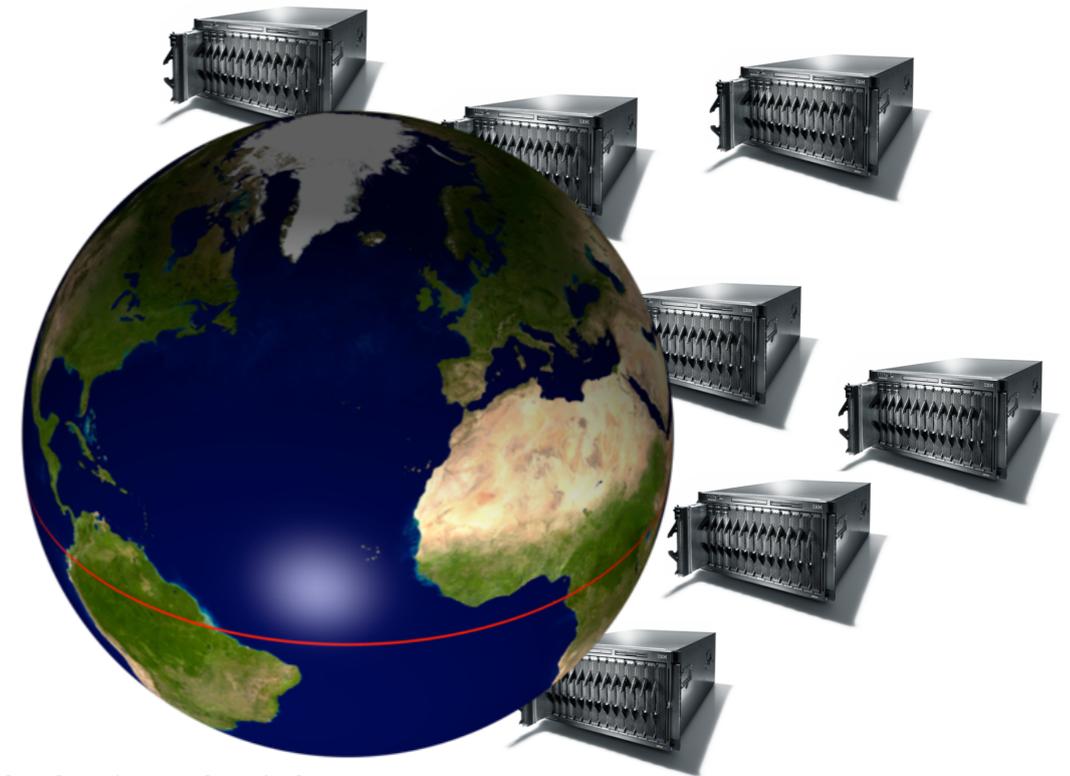
who m i

- Old Timer
- Early DNSSEC Reviewer
- NLnet Labs Guy

Overview

- DNS Overview
- DNSSEC in general
- The RR's
- Chain of Trust
- How to put it together
- Debugging aids

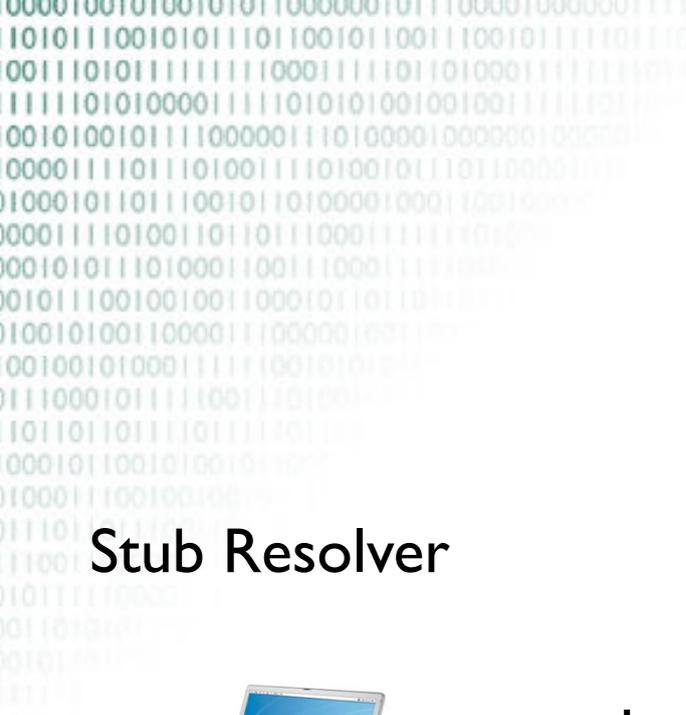
Authoritative Nameservers



Stub Resolver

Recursive Nameservers





Authoritative Nameservers **ROOT**



www.nlnetlabs.nl A

Recursive Nameservers



referral: nl NS

NL

Stub Resolver



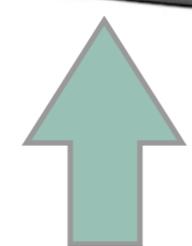
www.nlnetlabs.nl A

www.nlnetlabs.nl A



www.nlnetlabs.nl A 213.154.224.1

referral: nlnetlabs.nl NS



www.nlnetlabs.nl A 213.154.224.1

www.nlnetlabs.nl A

Answer: www.nlnetlabs.nl A 213.154.224.1



NLnetLabs.NL

root.hints: location of the root servers

```

; <<>> DiG 9.7.0b2 <<>> @k.root-servers.net www.nlnetlabs.nl
; (2 servers found)
;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 41886
;; flags: qr rd; QUERY: 1, ANSWER: 0, AUTHORITY: 7, ADDITIONAL: 12
;; WARNING: recursion requested but not available

;; QUESTION SECTION:
;www.nlnetlabs.nl. IN A

;; AUTHORITY SECTION:
nl.          172800 IN NS  nl1.dnsnode.net.
nl.          172800 IN NS  ns1.nic.nl.
nl.          172800 IN NS  ns2.nic.nl.
nl.          172800 IN NS  ns3.nic.nl.
nl.          172800 IN NS  ns4.nic.nl.
nl.          172800 IN NS  ns-nl.nic.fr.
nl.          172800 IN NS  sns-pb.isc.org.

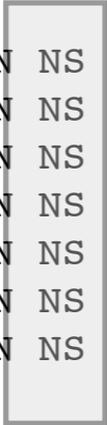
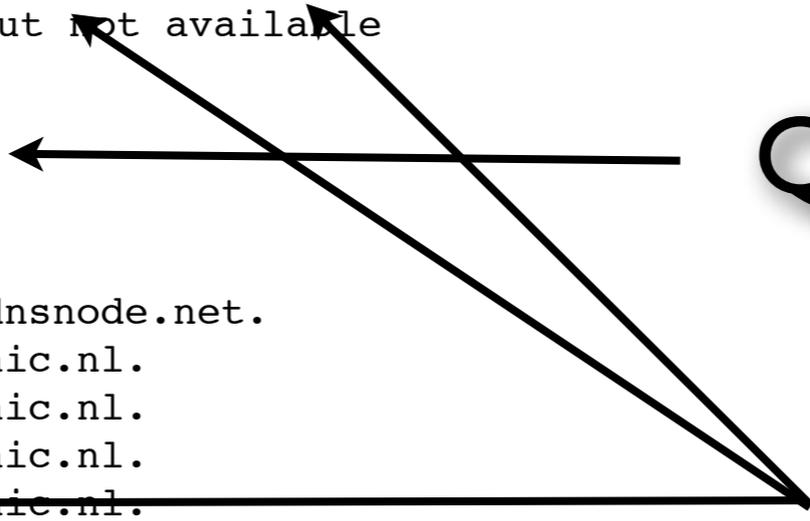
;; ADDITIONAL SECTION:
nl1.dnsnode.net. 172800 IN A 194.146.106.42
ns1.nic.nl.      172800 IN A 193.176.144.2
ns2.nic.nl.      172800 IN A 213.154.241.28
ns3.nic.nl.      172800 IN A 194.171.17.2
ns4.nic.nl.      172800 IN A 62.4.86.232
ns-nl.nic.fr.    172800 IN A 192.93.0.4
sns-pb.isc.org.  172800 IN A 192.5.4.1
ns1.nic.nl.      172800 IN AAAA 2a00:d78::102:193:176:144:2
ns2.nic.nl.      172800 IN AAAA 2001:7b8:606::28
ns3.nic.nl.      172800 IN AAAA 2001:610:0:800d::2
ns-nl.nic.fr.    172800 IN AAAA 2001:660:3005:1::1:2
sns-pb.isc.org.  172800 IN AAAA 2001:500:2e::1

;; Query time: 4 msec
;; SERVER: 2001:7fd::1#53(2001:7fd::1)
;; WHEN: Tue Apr 6 14:12:44 2010
;; MSG SIZE rcvd: 447

```

Question

Referral



Cache and TTL

```
;; ANSWER SECTION:  
www.nlnetlabs.nl. 10200 IN A 213.154.224.1
```

- TTL is a parameter that indicates how long data is to remain in a cache
- TTL value is set by the zone owner
- TTL decreases while in the cache

Cache Poison

- Attack is based on ‘predicting’ properties
 - e.g. when asking a question to a female you expect a female voice to answer
- If you ask a question with a specific QID you expect that QID in the answer
- Cache poisoner will take a wild guess

Isn't Query ID only sufficient?

Chance that n people have different birthdays

$$\bar{p}(n) = 1 \times \left(1 - \frac{1}{365}\right) \times \left(1 - \frac{2}{365}\right) \cdots \left(1 - \frac{n-1}{365}\right) = \frac{365 \times 364 \cdots (365 - n + 1)}{365^n} = \frac{365!}{365^n (365 - n)!}$$

Chance that n people have the same birthday

$$p(n) = 1 - \bar{p}(n).$$

n	P(n)
10	11,17%
20	41,1%
23	50,7%
30	70,6%
50	97%
57	99%
100	99,99997%

from: http://en.wikipedia.org/wiki/Birthday_problem

Bits	50%	5%	Aka
16	10 s	1 s	Unpatched server, random ID
26	2.8 h	17 m	Patched, using only 1024 ports
34	28 days	2.8 days	unbound with defaults
44	28444 days	2844.4 days	unbound with 0x20 and source addresses configured

Kaminsky's variant

- Classic cache poisoning gave you 'a few tries' to get in between the outgoing question and incoming answer
- Kaminsky came with a scheme where the culprit can keep trying
- Surprisingly simple, a wonder nobody thought of the variety before

DNSSEC

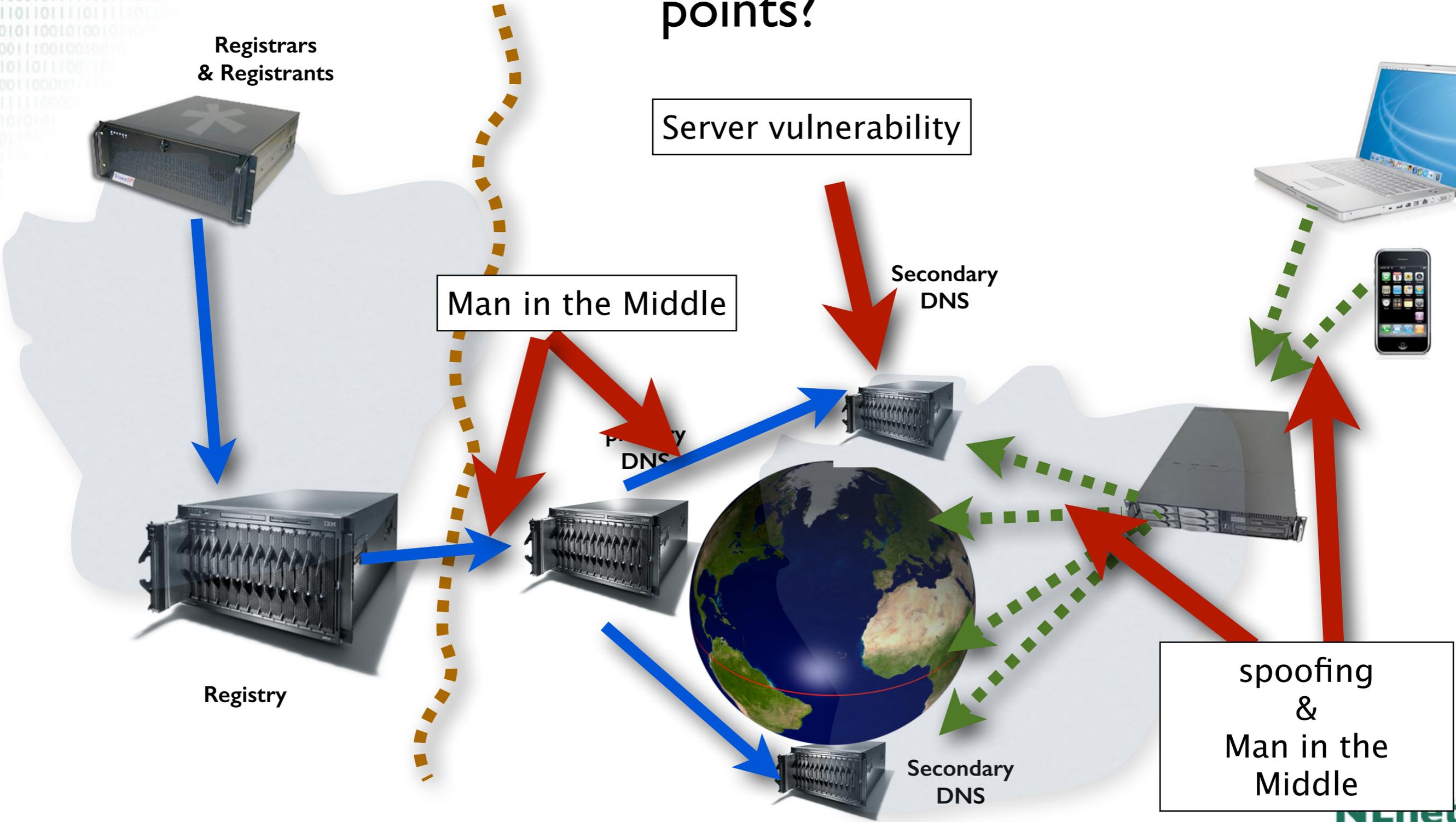
- Prevents cache poisoning for existing labels
- Proves non-existence of labels
- Attacks à la Kaminsky showed the need

DNSSEC Mechanisms

- New Resource Records
- Secure the Zone
- Delegating Signing Authority

Data flow through the DNS

Where are the vulnerable points?



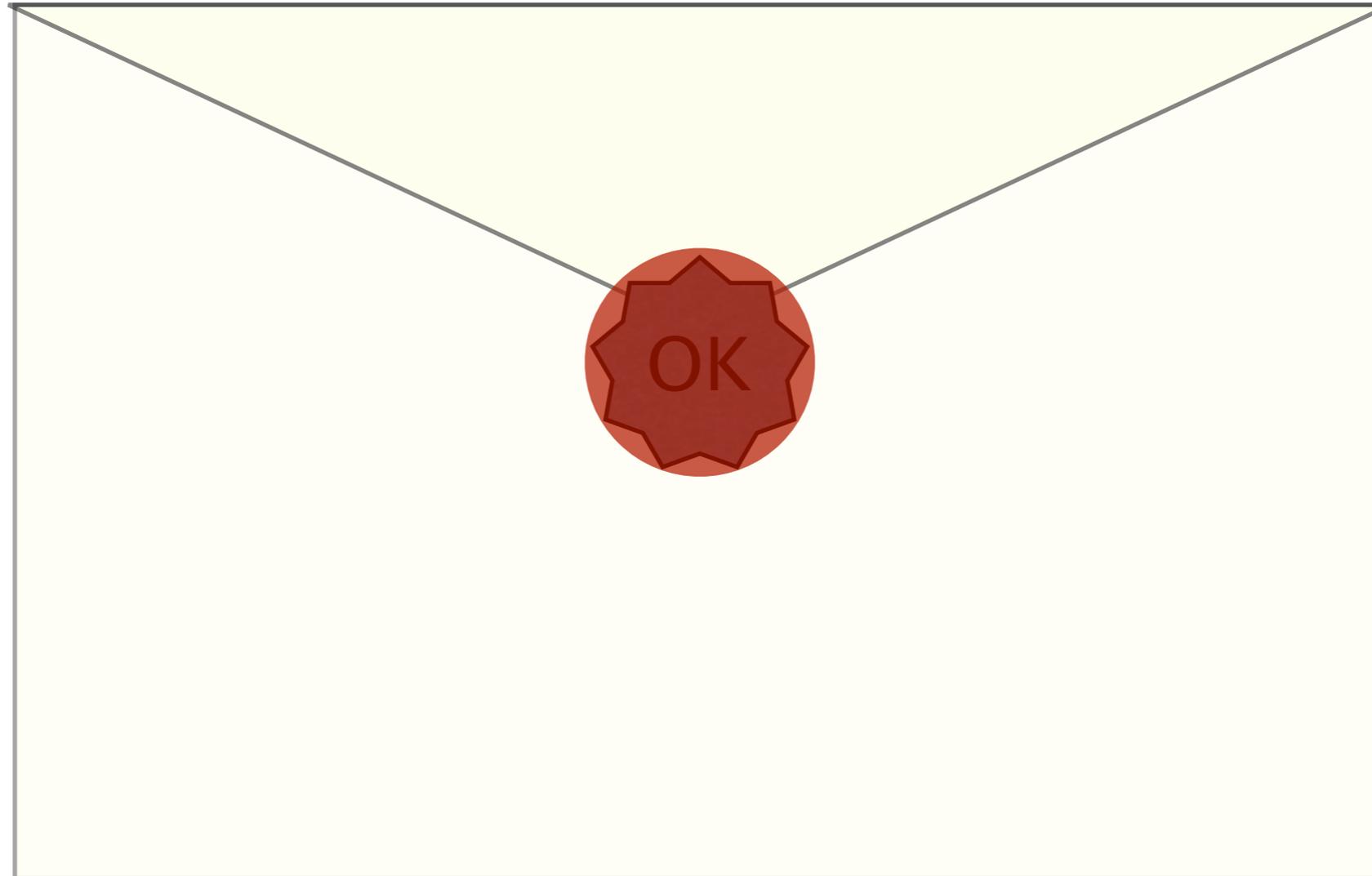
DNSSEC protects all this end-to-end

- As an aside:
There is a protection mechanism against the man in the middle: TSIG
- Provides hop-by-hop security
- TSIG is operationally deployed today
- Based on shared secret: not scalable
- Used a lot for AXFR transactions

What does DNSSEC provide

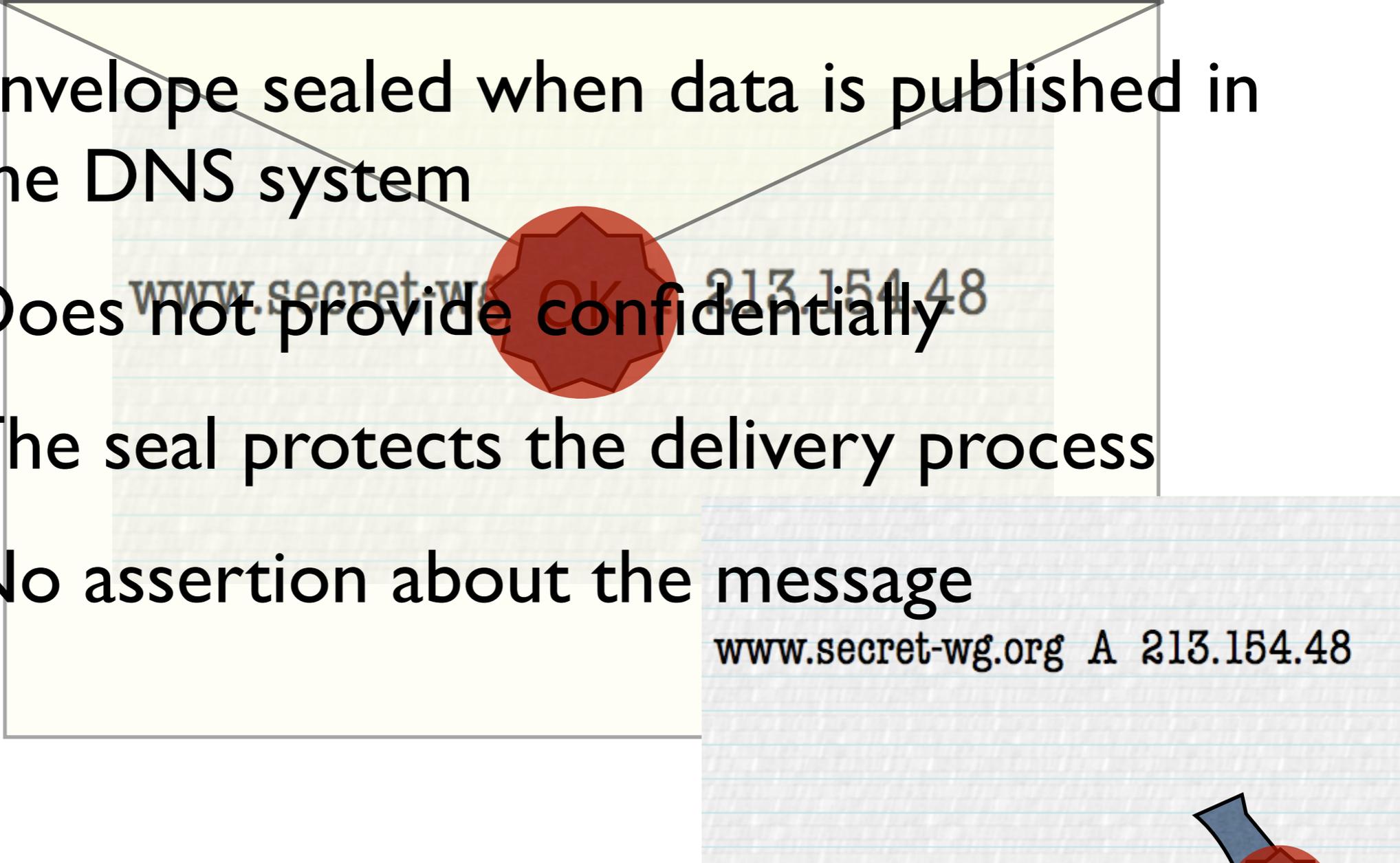
- Provides message authentication and integrity verification through cryptographic signatures
 - You know who provided the signature
 - No modifications between signing and validation
- It does not provide authorization
- It does not provide confidentiality
- It does not provide protection against DDOS

Metaphor



Metaphor

- Envelope sealed when data is published in the DNS system
- Does not provide confidentiality
- The seal protects the delivery process
- No assertion about the message



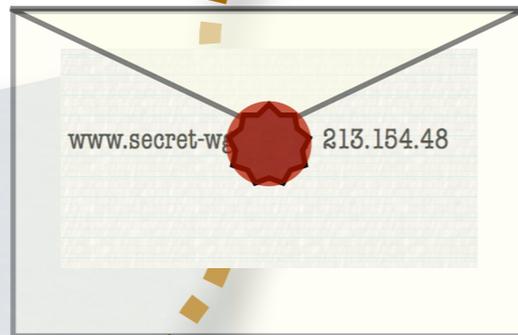
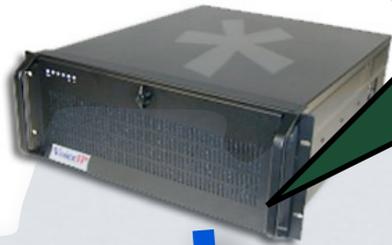
www.secret-wg.org A 213.154.48

Data flow through the DNS

End to end security

Apply Seal

Reg
& Registr



Registry



primary
DNS



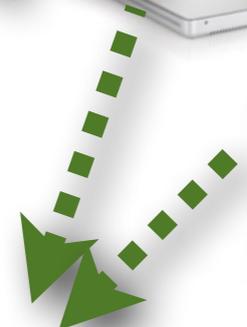
Secondary
DNS



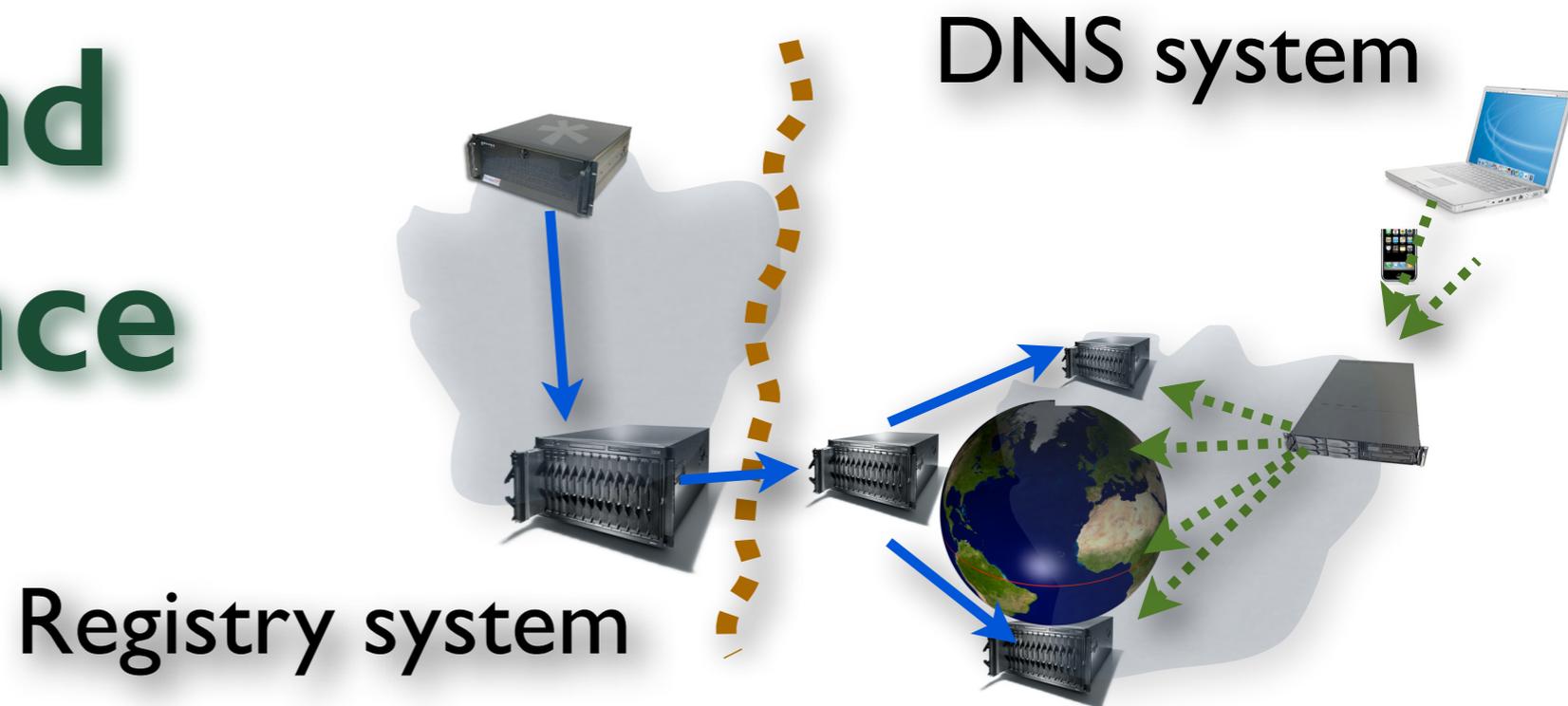
Validate Seal



Secondary
DNS



Trust and Confidence



- DNSSEC enables confidence in the DNS
- It does not change the trust we put in the Registry/Registrar procedures
- Although introduction of DNSSEC may improve some of the procedures

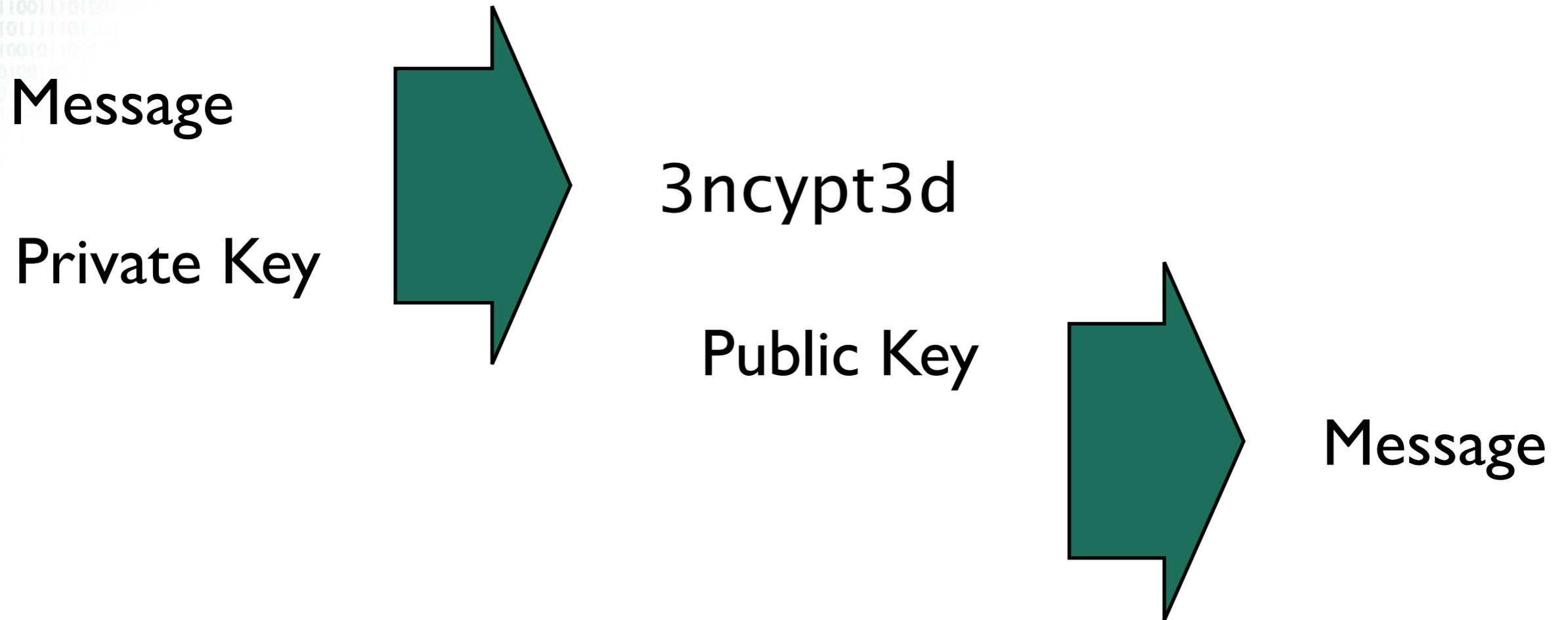
The mechanism used

- Using public key cryptographic algorithms signatures are applied over the DNS data
- By comparing the signatures with public keys the integrity and authenticity of the data can be established.

Public key cryptography in a nutshell

- Two large numbers and an encryption/decryption algorithm
- One of the numbers (the private key) and a message are used for encryption
- The other number (public key) and the decryption algorithm can be used to retrieve the original message

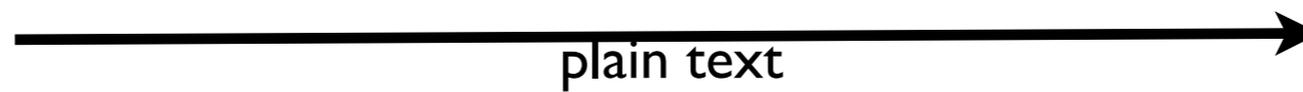
Public Crypto



Works only with matching key:
If you can decrypt with a public key you may
assert the message was signed with
corresponding private key

Use that method for signatures

Message



Message

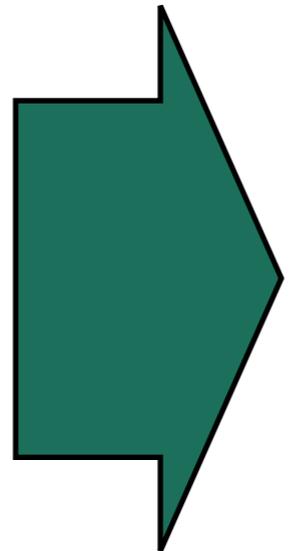


Calculated
Message
Digest

Decrypted
Message
Digest

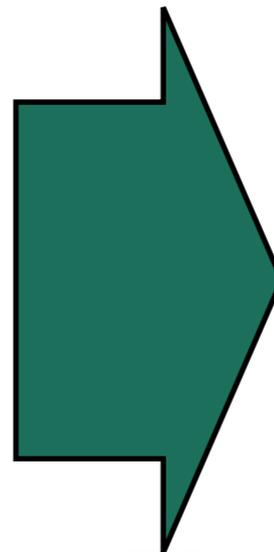


Message Digest



Private Key

Signature



Public Key

In Practice

- Key generation and signing is done by tools
- Validating and signing entity need to communicate which algorithms for hashing and public key cryptography is needed: e.g. RSASHA1, RSASHA256 or DSA

Holy Trinity

- Private Key: kept private and stored locally
- Public Keys: Published in the DNS as a DNSKEY Resource Record
- Signatures: Published in the DNS as a RRSIG Resource Record

Signing is done per Zone

- Each zone has one or more key-pairs for signing
- If you have the public keys from a zone you can validate signatures made with the corresponding private keys
- However, signing a complete zone does not scale

RRs and RRSets

- Resource Record:

```
— name          TTL    class  type  rdata
  www.nlnetlabs.nl.  7200   IN     A     192.168.10.3
```

- RRset: RRs with same name, class and type:

```
www.nlnetlabs.nl.  7200   IN     A     192.168.10.3
                  A     10.0.0.3
                  A     172.25.215.2
```

- RRsets are the atomic data units in the DNS
- RRsets are signed, not the individual RRs

DNSKEY RDATA

- 16 bits: FLAGS
- 8 bits: protocol
- 8 bits: algorithm
- $N \times 32$ bits: public key

`nlnetlabs.nl. 3600 IN DNSKEY 256 3 5 (`

```
AQ0vhvXXU61Pr8sCwELcqqq1g4JJ  
CALG4C9EtraBKVd+vGIF/unwigfLOA  
O3nHp/cgGrG6gJYe8OWKYNgq3kDChN)
```

RRSIG RDATA

- 16 bits - type covered
- 8 bits - algorithm
- 8 bits - nr. labels covered
- 32 bits - original TTL

```
nlnetlabs.nl. 3600 IN RRSIG A 5 2 3600 (
  20050611144523 20050511144523 3112 nlnetlabs.nl.
  VJ+8ijXvbrTLeoAiEk/qMrdudRnYZM1VlqhN
  vhYuAcYKe2X/jqYfMfjfSURmhPo+0/GOZjW
  66DJubZPmNSYXw== )
```

- 32 bit - signature expiration
- 32 bit - signature inception
- 16 bit - key tag
- signer's name

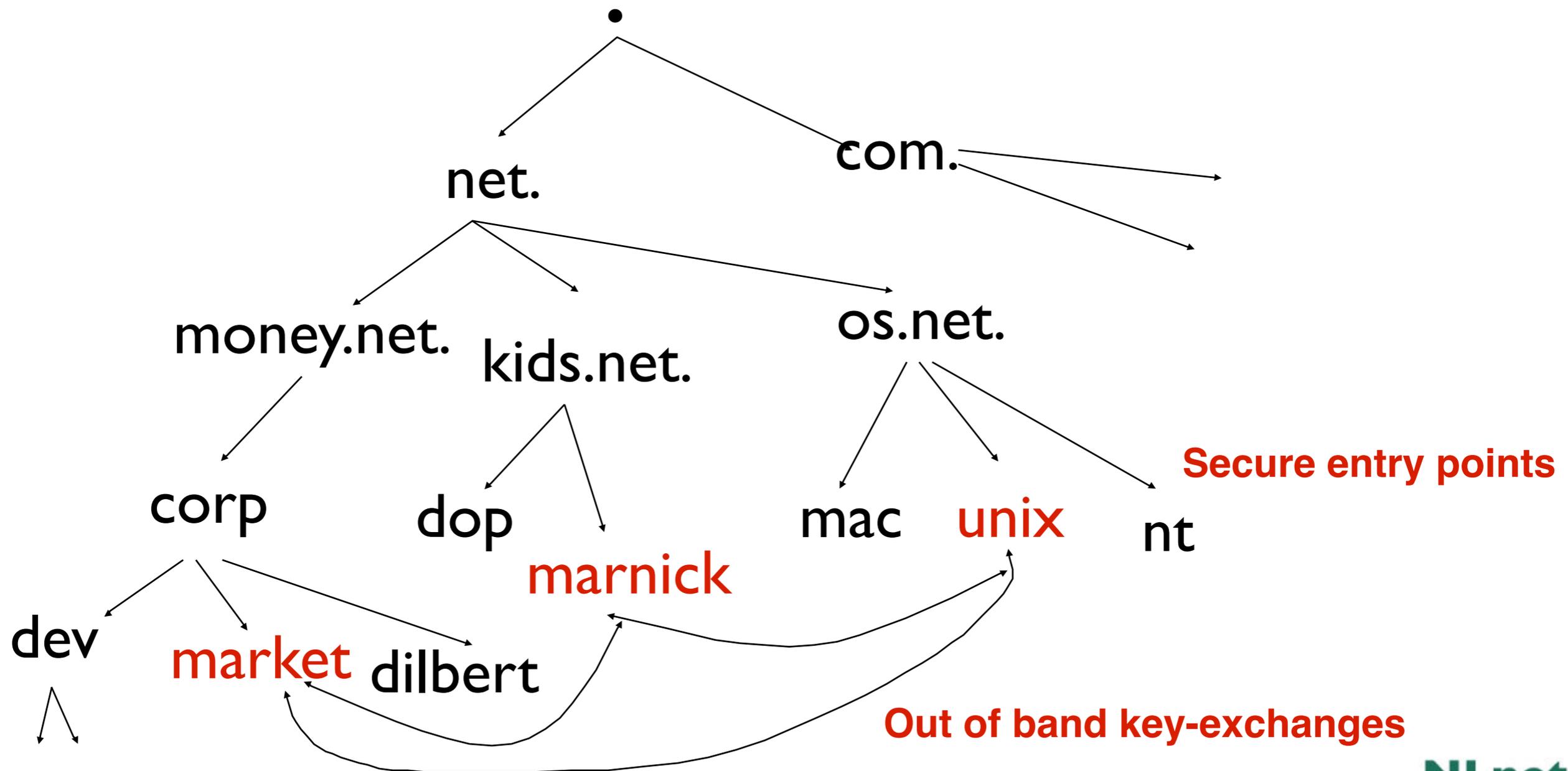
Validate Public Keys

- Make sure you get them from the appropriate entity and configure them as trust-anchors
- If you validate against the wrong public key there is a problem again
- For DNSSEC: key distribution through the DNS
 - Ideally only one key needed: that of the root of the DNS hierarchy (more on that later)

Delegating Signing Authority

Validating against configured keys

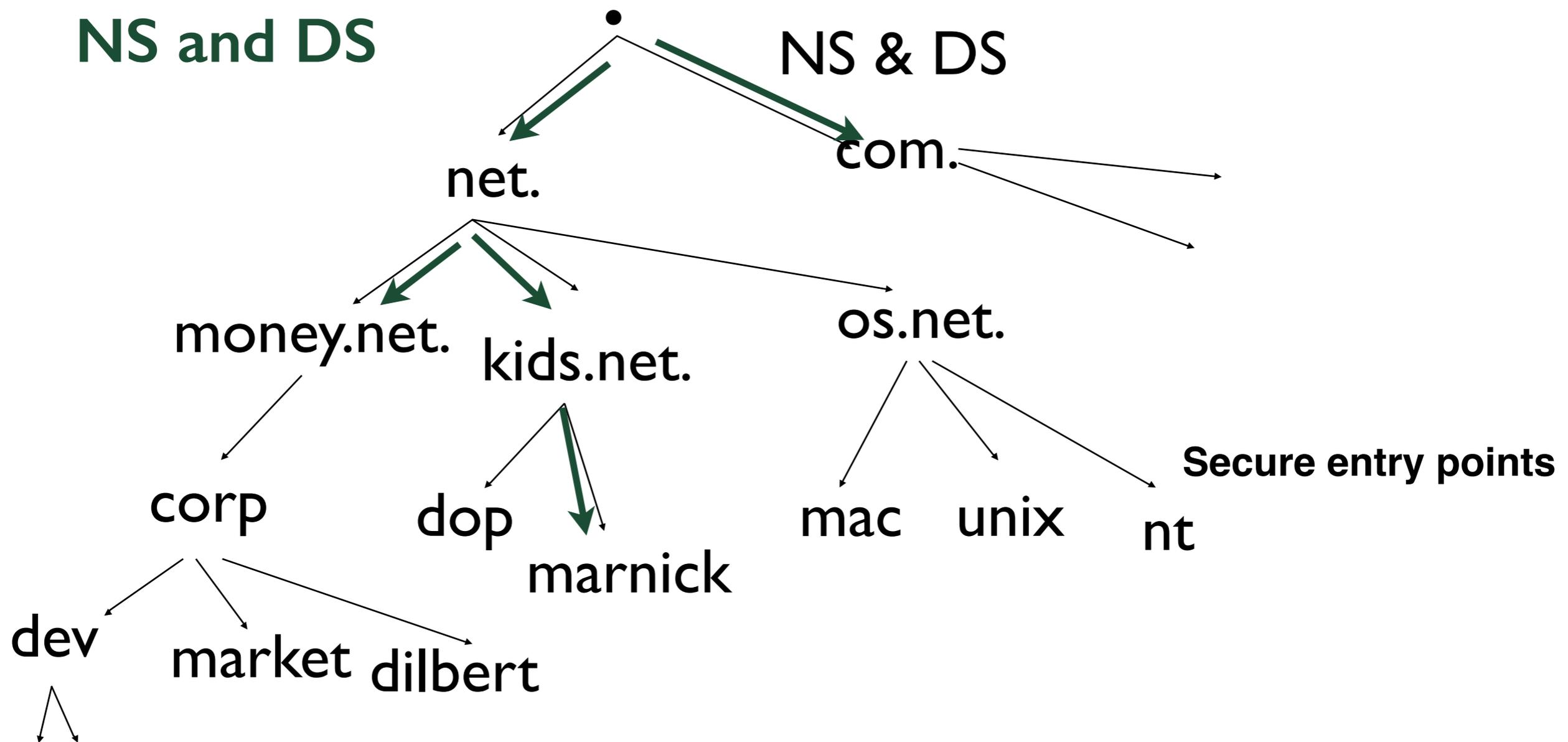
- Key distribution does not scale!



Locally Secured Zones

- Delegate Signing Security

NS and DS



Using the DNS to Distribute Keys

- Secured islands make key distribution problematic
- Distributing keys through DNS:
 - Use one trusted key to establish authenticity of other keys
 - Building chains of trust from the root down
 - Parents need to sign the keys of their children
- Only the root key needed in ideal world
 - Parents always delegate security to child

Delegation Signer (DS)

- Delegation Signer (DS) RR indicates that:
 - delegated zone is digitally signed
 - indicated key is used for the delegated zone
- Parent is authoritative for the DS of the child's zone
 - Not for the NS record delegating the child's zone!
 - DS should not be in the child's zone

DS RDATA

- 16 bits: key tag
- 8 bits: algorithm
- 8 bits: digest type
- 20 bytes: SHA-1 Digest

```
$ORIGIN nlnetlabs.nl.  
lab.nlnetlabs.nl. 3600 IN NS ns.lab.nlnetlabs.nl  
lab.nlnetlabs.nl. 3600 IN DS 3112 5 1 (  
239af98b923c023371b52  
1g23b92da12f42162b1a9 )
```

Key Problem

- Interaction with parent administratively expensive
 - Should only be done when needed
 - You might want to lock these in hardware
- Signing zones should be fast
 - Memory restrictions
 - Space and time concerns
 - Operational exposure higher

More Than One Key: KSK and ZSK

- RRsets are signed, not RRs
- DS points to specific key
 - Signature from that key over DNSKEY RRset transfers trust to all keys in DNSKEY RRset
- Key that DS points to only signs DNSKEY RRset
 - Key Signing Key (KSK)
- Other keys in DNSKEY RRset sign entire zone
 - Zone Signing Key (ZSK)

The Important Considerations

- KSK and ZSK have different 'shielding' properties:
KSK on smartcard, ZSK on disk
- ZSK needs 'daily' or permanent use.
- KSK less frequent
- ZSK change needs no involvement with 3rd parties
- KSK may need uncontrolled cooperation from 3rd parties

Initial Key Exchange

- Child needs to:
 - Send key signing keyset to parent
- Parent needs to:
 - Check child's zone
 - for DNSKEY & RRSIGs
 - Verify if key can be trusted
 - Generate DS RR

Walking the Chain of Trust

Locally configured
Trusted key: . 8907

\$ORIGIN .

1
2
3
4
5
6
7
8
9

```
. DNSKEY (...) 5TQ3s... (8907) ; KSK  
DNSKEY (...) lasE5... (2983) ; ZSK
```

```
RRSIG DNSKEY (...) 8907 . 69Hw9..
```

```
net. DS 7834 3 1ab15...  
RRSIG DS (...) . 2983
```

\$ORIGIN net.

```
net. DNSKEY (...) q3dEw... (7834) ; KSK  
DNSKEY (...) 5TQ3s... (5612) ; ZSK
```

```
RRSIG DNSKEY (...) 7834 net. cMas...
```

```
foo.net. DS 4252 3 1ab15...  
RRSIG DS (...) net. 5612
```

\$ORIGIN foo.net.

```
foo.net. DNSKEY (...) rwx002... (4252) ; KSK  
DNSKEY (...) sovP42... (1111) ; ZSK
```

```
RRSIG DNSKEY (...) 4252 foo.net. 5t...
```

```
www.foo.net. A 193.0.0.202  
RRSIG A (...) 1111 foo.net. a3...
```

Chain of Trust Verification, Summary

- Data in zone can be trusted if signed by a Zone-Signing-Key
- Zone-Signing-Keys can be trusted if signed by a Key-Signing-Key
- Key-Signing-Key can be trusted if pointed to by trusted DS record
- DS record can be trusted
 - if signed by the parents Zone-Signing-Key
 - or
 - DS or DNSKEY records can be trusted if exchanged out-of-band and locally stored (Secure entry point)

Where are we

- DNSKEY
- RRSIG
- DS

Offline Signing and Denial of Existence

- Problems with on-the-fly signing
 - Private key needs to be stored on an Internet facing system
 - Performance, signing is a CPU expensive operation
- How does one provide a proof that the answer to a question does not exist?

NSEC RDATA

- Points to the next domain name in the zone
 - also lists what are all the existing RRs for “name”
 - NSEC record for last name “wraps around” to first name in zone
- N*32 bit type bit map
- Used for authenticated denial-of-existence of data
 - authenticated non-existence of TYPEs and labels
- Example:

```
www.nlnetlabs.nl. 3600 IN NSEC nlnetlabs.nl. A RRSIG NSEC
```

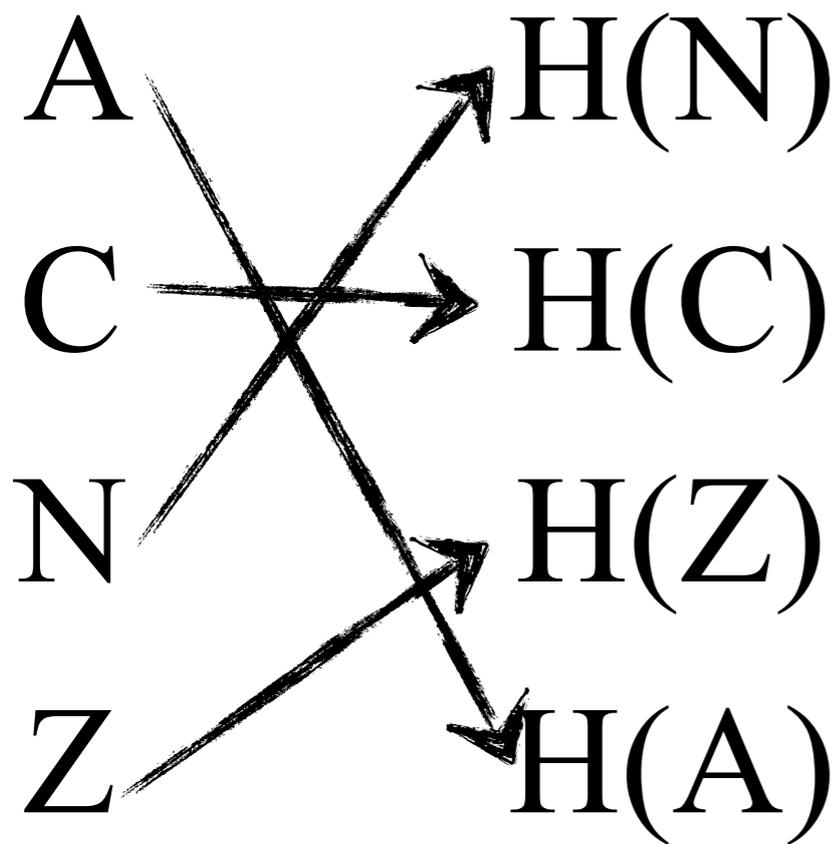
NSEC Records

- NSEC RR provides proof of non-existence
- If the servers response is Name Error (NXDOMAIN):
 - One or more NSEC RRs indicate that the name or a wildcard expansion does not exist
- If the servers response is NOERROR:
 - And empty answer section
 - The NSEC proves that the QTYPE did not exist
- More than one NSEC may be required in response
 - Wildcards
- NSEC records are generated by tools
 - Tools also order the zone

NSEC Walking

- NSEC records allow for zone enumeration
 - Providing privacy was not a requirement at the time
 - Zone enumeration is a deployment barrier
-
- Solution has been developed: NSEC3
 - RFC 5155
 - Complicated piece of protocol work
 - Hard to troubleshoot
 - Only to be used over Delegation Centric Zones

NSEC3



- Creates a linked list of the hashed names
- Non-existence proof of the hash proofs non-existence of original
- Dictionary attack barriers:
 - Salt
 - Iterations

New Resource Records

- Three Public key crypto related RRs
 - RRSIG: Signature over RRset made using private key
 - DNSKEY: Public key, needed for verifying a RRSIG
 - DS: Delegation Signer; 'Pointer' for building chains of authentication

- One RR for internal consistency
 - NSEC and NSEC3: Indicates which name is the next one in the zone and which typecodes are available for the current name
 - authenticated non-existence of data

Other Keys in the DNS

- DNSKEY RR can only be used for DNSSEC
 - Keys for other applications need to use other RR types
- CERT
 - For X.509 certificates
- Application keys under discussion/development
 - IPSECKEY
 - SSHFP Summary for now
 - DANE!!!

Practicalities

- Create the keys
- Sign the zones
- Keep everything up to date

Key creation

```
$ ldns-keygen -a RSASHA256 -b 1024 1sand0s.nl
K1sand0s.nl.+008+24201
$ ls
K1sand0s.nl.+008+24201.ds K1sand0s.nl.+008+24201.private
K1sand0s.nl.+008+24201.key
$ ldns-keygen -k -a RSASHA256 -b 1024 1sand0s.nl
K1sand0s.nl.+008+24040
$ ls
K1sand0s.nl.+008+24040.ds K1sand0s.nl.+008+24201.ds
K1sand0s.nl.+008+24040.key K1sand0s.nl.+008+24201.key
K1sand0s.nl.+008+24040.private
K1sand0s.nl.+008+24201.private
```

Zone signing

```
$ ldns-signzone 1sand0s.nl.zone K1sand0s.nl.+008+24040 \  
                K1sand0s.nl.+008+24201  
$ ls  
1sand0s.nl.zone K1sand0s.nl.+008+24040.private  
1sand0s.nl.zone.signed K1sand0s.nl.+008+24201.ds  
K1sand0s.nl.+008+24040.ds k1sand0s.nl.+008+24201.key
```

Notify Parent

Sending the DS to the Parent to attach to the chain of trust

Parent puts this in his/her zone

Signs the zone and publish etc ...

Resigning needs

- When zone changes
- New contents
- When RRSIG expires
- Idns-signzone 30 days
- New chore for sys admin!
 - Automate
 - (ask for a raise)

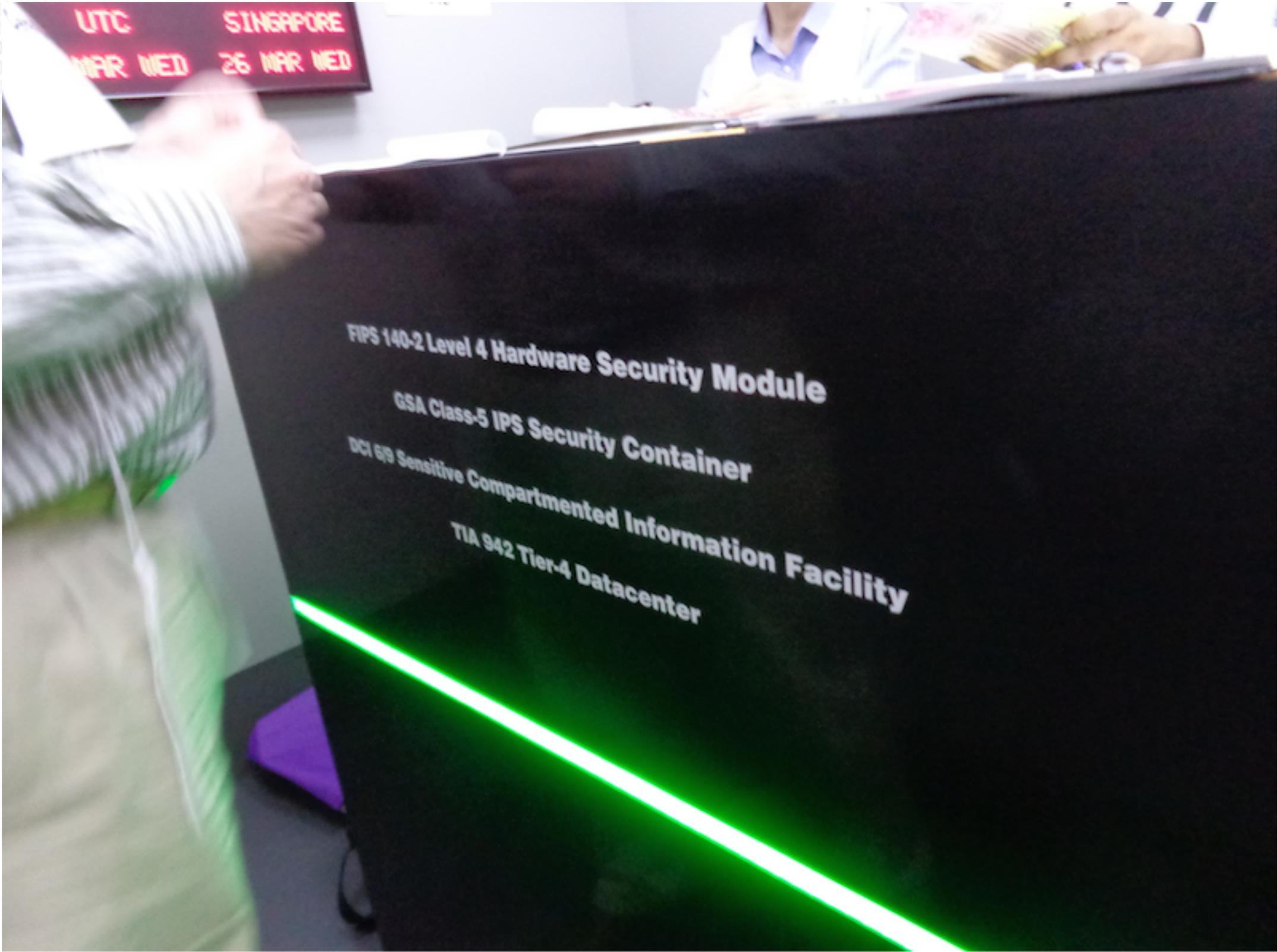
Key rollovers

DNSKEY flavours

- Zone Signing Key (ZSK)
- Key Signing Key (KSK)
- Functions as secure entry point into the zone
 - Trust-anchor configuration
 - Parental DS points to it
 - Interaction with 3rd party/parties
- DNSKEYs are treated all the same in the protocol
- Operators can make a distinction
 - Look at the flag field: ODD (257 in practice) means SEP

Benefits of using separate keys

- Rolling KSK needs interaction, rolling ZSKs can be done almost instantaneously
- Remember KSK replacement may result in
 - Trust-anchor updates
 - Change of DS record at parent
- Allows different responsibilities
 - ZSKs may be touched day to day by junior staff
 - KSKs may only be touched by senior staff









Rolling keys instantaneously?

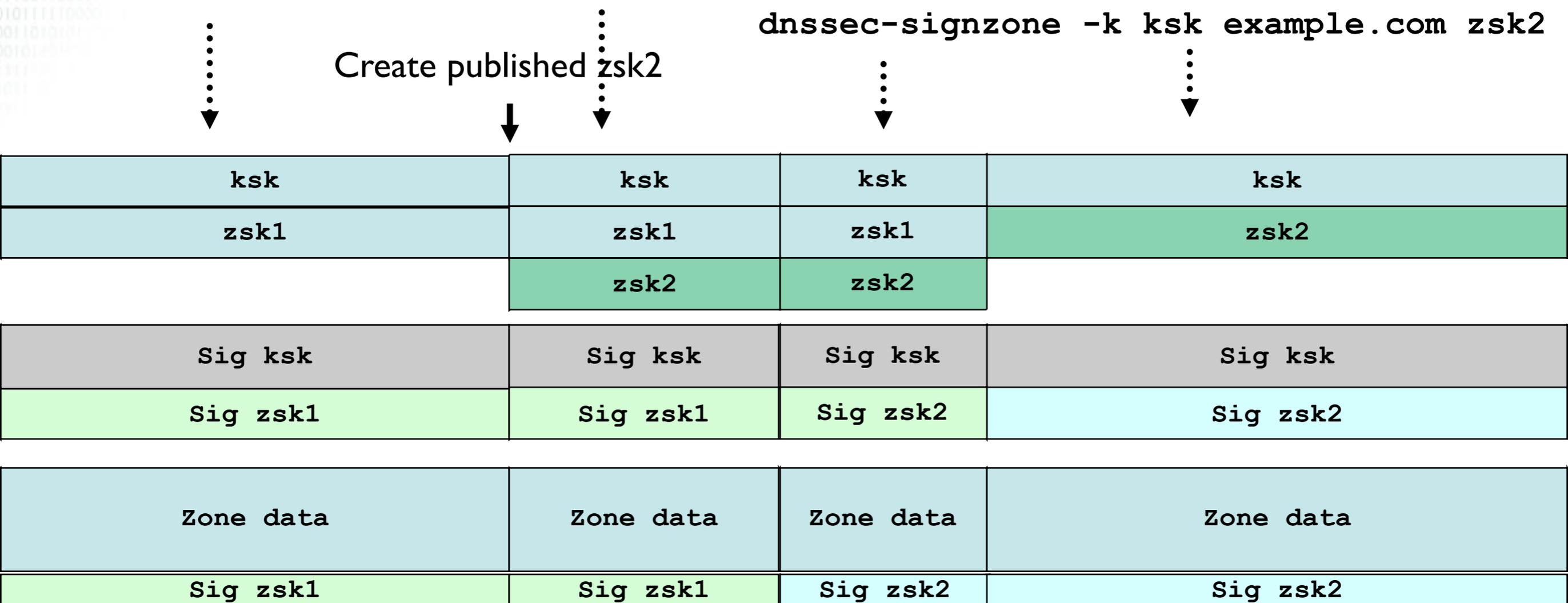
- Remember that in the DNS caches are at play.
- It takes a bit of time to have new information propagate
- When you happen to get new DNSKEYs you would like to be able to use DNSSIGs from the cache
- When you happen to get old DNSKEYs from the cache you would like to use new DNSSIGs
- Try to make sure both old and new keys are available
- Or, try to make sure both old and new sigs are available

ZSK rollover

```
dnssec-signzone -k ksk example.com zsk1
```

```
dnssec-signzone -k ksk example.com zsk2
```

Create published zsk2



time

At least TTL of DNSKEY RRs

At least MAX TTL over all RRs

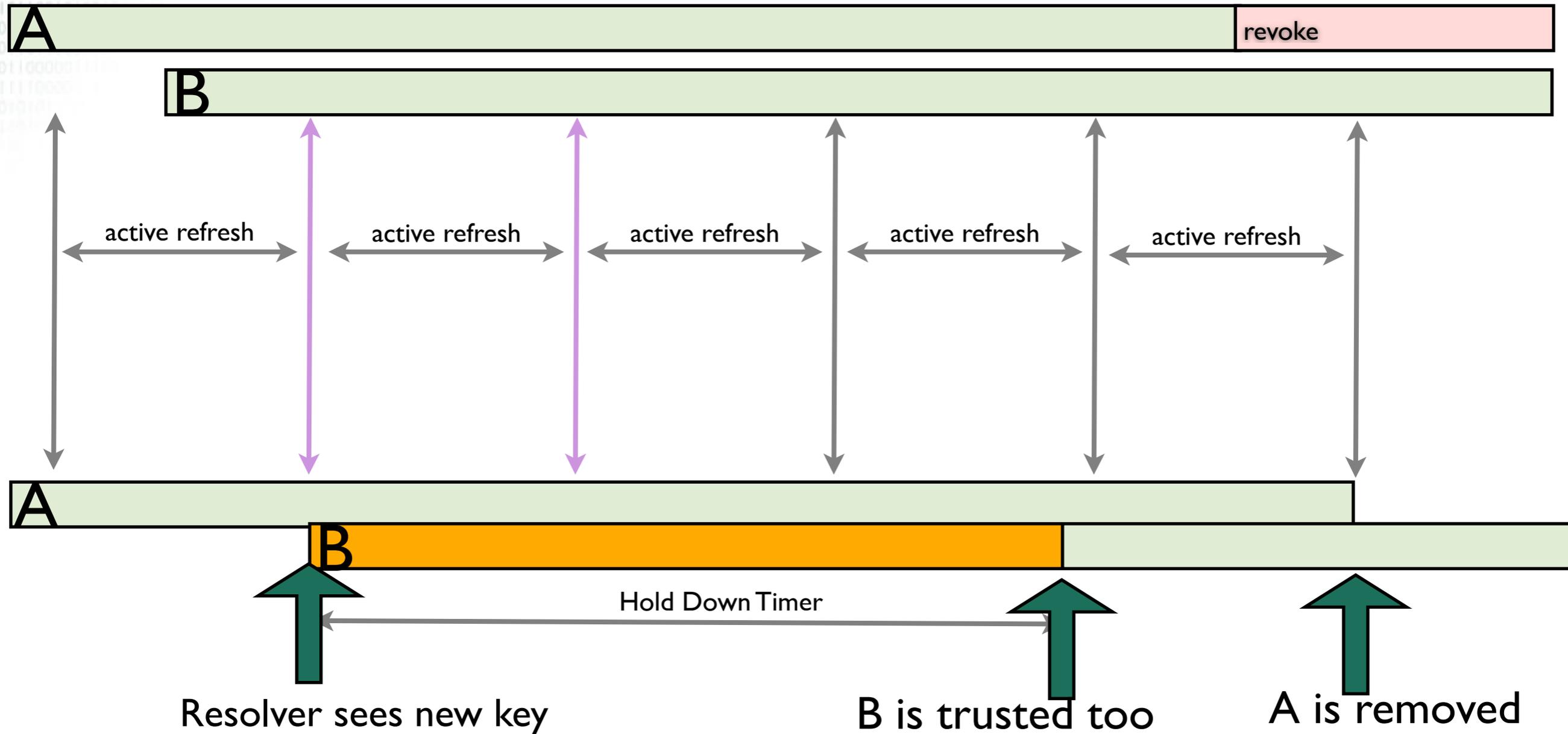
KSK rollover

- You are dependent on your parent.
- You cannot control when the parent changes the DS rr
- Use the old KSK until the old DNS had time to propagate from caches

RFC 5011 Concepts

- Trust anchor maintenance based on existing trust relation
- New keys only accepted after its been seen for more than 30 days (Hold Down)
- Signaling retirement of the key by setting a 'revoke' flag

Zone Owner



Trust Anchor state

Keeping up to date

By Hand?

- Error prone and complex
- Easy to forget

Use Tools!

- Cronjobs + Scripts
- Tools like OpenDNSSEC
- Build-in tools

Problem Shooting

Help, the Internet just stopped working

- DNS is the very resilient
- Install and forget
 - 90 % works by accident
- DNSSEC makes the (this) system brittle
- Only three results
 - Secure, insecure, “bogus”
 - Needs periodic maintenance

Most seen problems

- Expired SIGS
- Expired KEYS
- Missing DS (signed but indeterminable state)
- Faulty clocks
- Algorithms mismatch
 - Incomplete rollovers

Debugging tools: dig, date, pencil & paper

- Dig for the Records
- Check the dates
- Draw the chain and see where it brakes

NLnet Labs Tools

- Signing tools etc
- Drill: like dig but with DNSSEC integrated
- Bottom up (-S) or Top down (-TD)
DNSSEC chain checking
- Unbound-host -f <key> -d

Online Tools

- <http://dnssec-debugger.verisignlabs.com>
- <http://dnsviz.net>
- <http://dnscheck.iis.se> (<http://dnscheck.se>)
- <http://www.zonecheck.fr/demo/>

More Info

<http://www.internetsociety.org/deploy360/dnssec/>

<http://workbench.sidnlabs.nl>

<https://dnssec.surfnet.nl>



Questions

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